

– Workshop on Stochastics and Special Functions –

Combinatorics of Orthogonal Polynomials

– Perturbed Chebyshev polynomials –

Richard Brak

Mathematics and Statistics of Complex systems
Department of Mathematics and Statistics
University of Melbourne

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Orthogonal Polynomials: Two questions

- Definition (Moment functional)

Let $\{\mu_n\}_{n \geq 0}$, $\mu_n \in \mathbb{C}$ be a sequence, called the *moment sequence*, and $\mathcal{L} : \mathbb{C}(x) \rightarrow \mathbb{C}$ a functional acting on polynomials $\mathbb{C}(x)$. Then \mathcal{L} is called a moment functional if

- 1 \mathcal{L} is linear
- 2 $\mathcal{L}(x^n) = \mu_n$, $n \geq 0$

- Definition (Orthogonal Polynomials)

Let $\{P_k(x)\}_{k \geq 0}$ be a sequence of polynomials in x with coefficients in the field \mathbb{C} . The polynomials are orthogonal with respect to a moment functional \mathcal{L} if, for all $n, m \geq 0$,

- 1 $\deg(P_n) = n$
- 2 $\mathcal{L}(P_n P_m) = \Lambda_n \delta_{n,m}$, $\Lambda_n \in \mathbb{C} \setminus 0$

- Question: Given a moment sequence $\{\mu_n\}_{n \geq 0}$ does there exist an orthogonal polynomial sequence?
- Answer: Yes, so long as the Hankel determinants do not vanish.

$$H_n = \begin{pmatrix} \mu_0 & \cdots & \mu_n \\ \vdots & \ddots & \vdots \\ \mu_n & \cdots & \mu_{2n} \end{pmatrix}$$

then

$$P_n(x) = \frac{1}{\det H_{n-1}} \det \begin{pmatrix} \mu_0 & \cdots & \mu_n \\ \vdots & \ddots & \vdots \\ \mu_{n-1} & \cdots & \mu_{2n-1} \\ 1 & \cdots & x^n \end{pmatrix}.$$

- Question: Give a sequence of orthogonal polynomials $\{P_k(x)\}_{k \geq 0}$ that satisfy the classical three term recurrence does there exist a moment sequence and linear functional which they are orthogonal with respect to?

- Theorem (Favard)

Let $\{b_k\}_{k \geq 0}$ and $\{\lambda_k\}_{k \geq 1}$ be sequences with $b_k, \lambda_k \in \mathbb{C}$ and let $\{P_k(x)\}_{k \geq 0}$ be a sequence of polynomials satisfying

$$P_{k+1}(x) = (x - b_k)P_k(x) - \lambda_k P_{k-1}(x) \quad k \geq 1$$

$$P_0 = 1, \quad P_1 = x - b_0$$

Then there exists a unique moment functional $\mathcal{L} : x^n \rightarrow \mu_n$ such that $\{P_k(x)\}_{k \geq 0}$ are a monic orthogonal sequence with respect to \mathcal{L} iff $\lambda_n \neq 0$ for all $n \geq 1$.

- Compute $\{\mu_n\}_{n \geq 0}$ recursively $\mathcal{L}(1) = \mu_0$, $\mathcal{L}(P_n) = 0$, $n > 0$.

Orthogonal Polynomials: Combinatorics

- We would like a combinatorial representation:
 - Combinatorial computation of polynomials and moments
 - Combinatorial proofs; Favard, orthogonality etc.
 - Orthogonal polynomial results give combinatorial answers
- Two fundamental representations

$$P_n(x) = \sum_{\pi \in \mathbb{P}_n} w(\pi)$$

$$\mu_n = \sum_{\tau \in \mathbb{M}_n(0,0)} w(\tau)$$

\mathbb{P}_n = Set of pavings on a line of n vertices

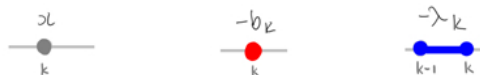
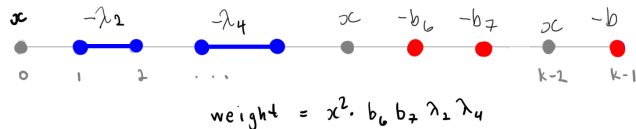
$\mathbb{M}_n(0,0)$ = Set of $0 \rightarrow 0$ Motzkin paths with n steps

and $w(\pi)$, $w(\tau)$ is the respective weights.

Combinatorial Objects

Definitions of pavings and paths – by example

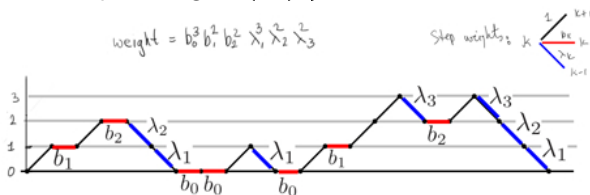
- Paving: Line segment with n vertices by:
 - 'monomers' – weights $\{b_k\}_{k \geq 0}$
 - 'dimers' – weights $\{\lambda_k\}_{k \geq 0}$



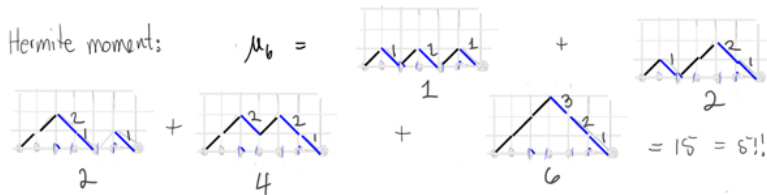
- Example: Hermite polynomials: $b_k = 0$, $\lambda_k = k$

$$k=3: \quad \begin{array}{c} x \\ \bullet \end{array} \begin{array}{c} x \\ \bullet \end{array} \begin{array}{c} x \\ \bullet \end{array} + \begin{array}{c} -1 \\ \bullet \end{array} \begin{array}{c} x \\ \bullet \end{array} + \begin{array}{c} x \\ \bullet \end{array} \begin{array}{c} -2 \\ \bullet \end{array} \Rightarrow H_3 = x^3 - 3x$$

- Motzkin paths: n steps: 'up', 'down' and 'horizontal'
 - down steps – weights $\{b_k\}_{k \geq 0}$
 - horizontal steps – weights $\{\lambda_k\}_{k \geq 0}$



- Example: Hermite polynomials: $b_k = 0$, $\lambda_k = k$



In general $\mu_n = (n-1)!!$ n even i.e. number of fixed point free involutions – bijection.

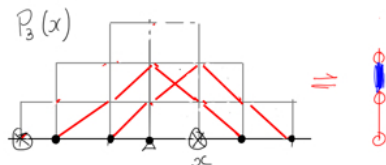
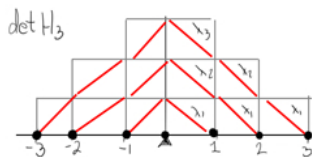
Dual representation

Hankel determinants

$$H_n = \begin{pmatrix} \mu_0 & \cdots & \mu_n \\ \vdots & \ddots & \vdots \\ \mu_n & \cdots & \mu_{2n} \end{pmatrix} \quad P_n(x) \sim \det \begin{pmatrix} \mu_0 & \cdots & \mu_n \\ \vdots & \ddots & \vdots \\ \mu_{n-1} & \cdots & \mu_{2n-1} \\ 1 & \cdots & x^n \end{pmatrix}.$$

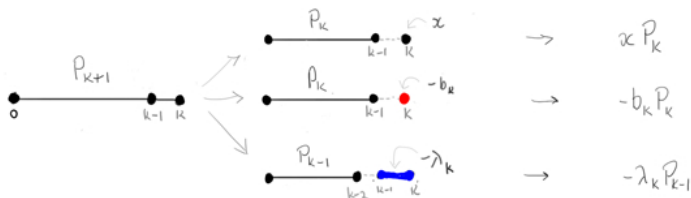
Combinatorially the KMLGV theorem gives:

$\det H_n =$ number of n 'non-intersecting' Motzkin paths from $-k$ to k .



Combinatorial proofs

- $P_{k+1}(x) = (x - b_k)P_k(x) - \lambda_k P_{k-1}(x)$
- $P_k(x) = \sum_{\pi \in \mathbb{P}_n} w(\pi)$



Continued fraction connection

- Generating function for moments ie. Motzkin paths:

$$M(t) = \sum_{n \geq 0} \mu_n t^n.$$

- Jacobi continued fractions:

$$M(t) = \frac{1}{1 - b_0 t - \frac{\lambda_1 t^2}{1 - b_1 t - \frac{\lambda_2 t^2}{\dots}}}$$

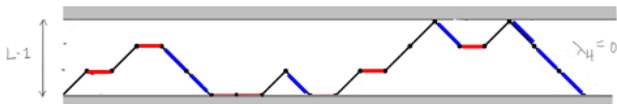
- Convergents – truncate at level L : set $\lambda_L \rightarrow 0$ and simplify

$$M(t) = \frac{N_{L-1}(t)}{D_L(t)}, \quad N_k(t) = \hat{P}_k^*(t), \quad D_n(t) = \hat{P}_n(t)$$

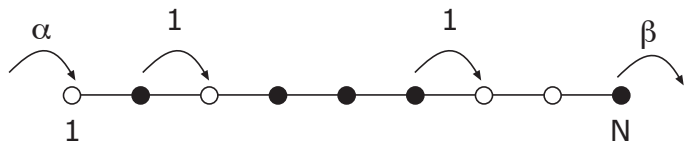
\hat{P}_n is the reciprocal orthogonal polynomial: $\hat{P}_n(t) = t^{-n} P_n(1/t)$

Setting $\lambda_L \rightarrow 0$ means Motzkin paths cannot have any steps at or above height L .

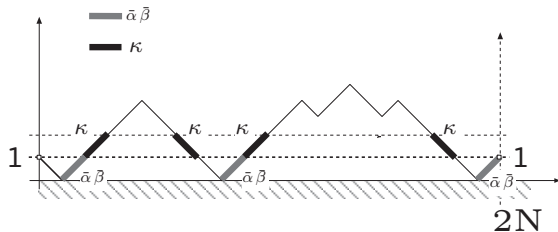
- $M(t) = \frac{N_{L-1}(t)}{D_L(t)}$ generates Motzkin paths in strip of width n



Asymmetric Simple Exclusion Process (ASEP).



- Stationary State normalisation: $Z_N(\alpha, \beta)$.

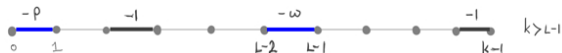


$\bar{\alpha} = 1/\alpha$, $\bar{\beta} = 1/\beta$, $\kappa^2 = 1 - (1 - \bar{\alpha})(1 - \bar{\beta})$:
 Motzkin with $\lambda_1 = \bar{\alpha}\bar{\beta}$ and $\lambda_2 = \kappa^2$, $\lambda_k = 1$, $k > 2$.

Perturbed Chebyshev

More generally, other problems (polymers interacting with a strip) require $b_k = 0$ and $\lambda_1 = -\rho$ and $\lambda_L = -\omega$, $\lambda_k = -1$, $k \neq 1, L$.
(joint work with J-A. Osborn)

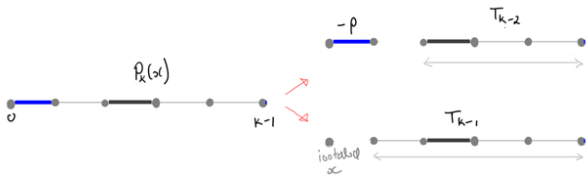
- If $\lambda_k = 1$, $k \geq 1$ then Chebyshev of the first kind.
- Thus we have the following paving problem to solve.



- But first $k < L - 1$



- Set of pavings partitions into two:



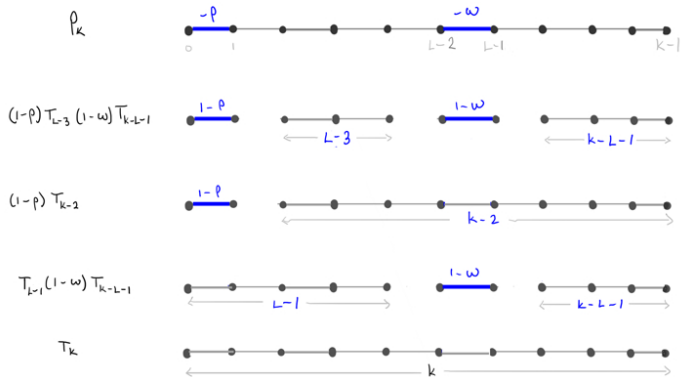
$$\begin{aligned}
 P_k &= (-\rho)T_{k-2} + xT_{k-1} \\
 &= (1-\rho)T_{k-2} + T_k
 \end{aligned}$$

using $xT_{k-1} = T_k + T_{k-2}$

- or combinatorially



- Thus for the two 'defect' case $k > L - 1$. There are four cases:



Thus we have,

Proposition

The orthogonal polynomial sequence P_n with $\lambda_1 = \rho$, $\lambda_{L-1} = \omega$ and $\lambda_k = 1$, $k \neq 1, L - 1$ is given by

$$P_k = T_k + (1 - \rho)T_{k-2} + (1 - \omega)T_{L-1}T_{k-L-1} \\ + (1 - \rho)(1 - \omega)T_{L-3}T_{k-L-1}$$

where $T_k(x)$ are the Chebyshev (like) polynomials

$$T_{k+1} = xT_k(x) - T_{k-1}(x), \quad P_1 = x, \quad P_0 = 1.$$

Thank You.